

17-423/723: Software System Design

Designing Interface Specifications

Jan 28, 2026



Logistics

- HW1 due today
- Project teams announced
- M1 released later today

Learning Goals

- Describe the role and importance of an interface specification
- Describe the structure and meaning of a specification
- Describe different dimensions to consider while designing a specification

Examples & figures based on <https://ocw.mit.edu/ans7870/6/6.005/s16/>

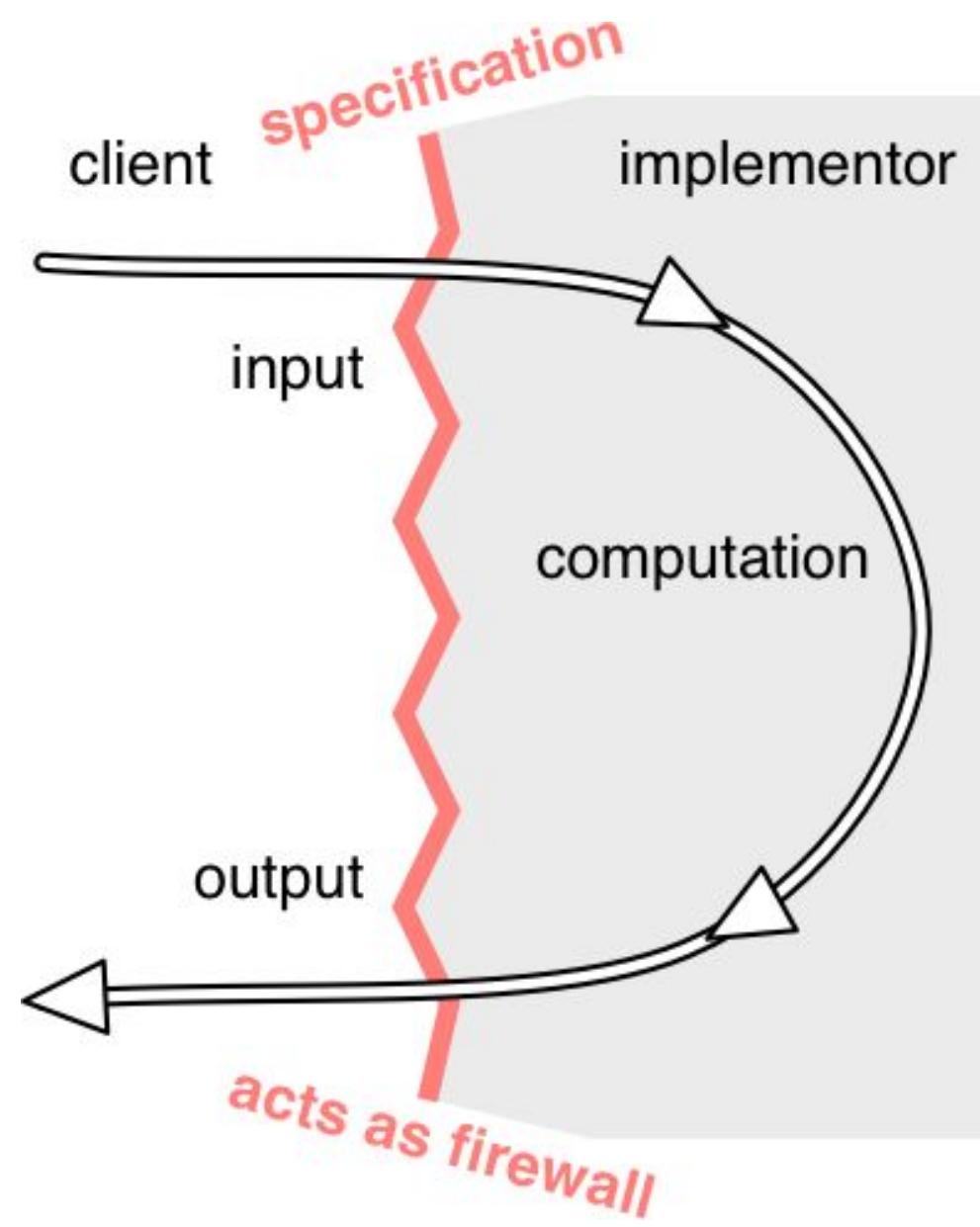
Interface Specifications

Specification

- A statement of a desired behavior or quality attribute of a software system
- Functional specification
 - “The scheduling system must provide a way for the patient to modify an existing appointment”
- Quality attribute specification
 - “The system must be able to handle additional 5000 users without a loss of latency” (scalability)
- **Interface specification**
 - Describes how a component interacts with its clients through one or more services that it provides
 - **Today's focus!**

Interface Specification

- **Contract** between a client and a component
- **For clients:**
 - Describes what a client needs to know to use the component
 - Describes what is expected as the output, given an input
 - Hides implementation details
- **For implementors:**
 - Describes implementation tasks to be fulfilled by developers (or LLMs)
 - Hides possible uses of the component by clients



Interface Specifications in Practice

OVERVIEW PACKAGE CLASS USE TREE DEPRECATED INDEX HELP

PREV CLASS NEXT CLASS FRAMES NO FRAMES ALL CLASSES

SUMMARY: NESTED | FIELD | CONSTR | METHOD DETAIL: FIELD | CONSTR | METHOD

compact1, compact2, compact3
java.util

Class HashSet<E>

java.lang.Object
java.util.AbstractCollection<E>
java.util.AbstractSet<E>
java.util.HashSet<E>

Method Summary

All Methods

Instance Methods

Concrete Methods

Modifier and Type

boolean

void

Object

boolean

boolean

Iterator<E>

Java Collections API

Method and Description

`add(E e)`

Adds the specified element to this set if it is not already present.

`clear()`

Removes all of the elements from this set.

`clone()`

Returns a shallow copy of this HashSet instance: the elements themselves

`contains(Object o)`

Returns true if this set contains the specified element.

`isEmpty()`

Returns true if this set contains no elements.

`iterator()`

Returns an iterator over the elements in this set.

Interface Specifications in Practice

```
def add(num1, num2):  
    """  
    Add up two integer numbers.  
  
    This function simply wraps the ``+`` operator, and does not  
    do anything interesting, except for illustrating what  
    the docstring of a very simple function looks like.  
  
    Parameters  
    -----  
    num1 : int  
        First number to add.  
    num2 : int  
        Second number to add.  
  
    Returns  
    -----  
    int  
        The sum of ``num1`` and ``num2``.  
  
    See Also  
    -----  
    subtract : Subtract one integer from another.  
  
    Examples  
    -----  
    >>> add(2, 2)  
    4  
    >>> add(25, 0)  
    25  
    >>> add(10, -10)  
    0  
    """
```

Python Docstrings

Interface Specifications in Practice

Swagger Petstore

1.0.7 OAS 2.0

[Base URL: petstore.swagger.io/v2]
<https://petstore.swagger.io/v2/swagger.json>

This is a sample server Petstore server. You can find out more about Swagger

store Access to Petstore orders

GET /store/inventory Returns pet inventories by status

POST /store/order Place an order for a pet

Parameters

Name	Description
body * required	order placed for purchasing the pet
object (body)	Example Value Model

```
{  
  "id": 0,  
  "petId": 0,  
  "quantity": 0,  
  "shipDate": "2025-02-03T21:44:39.405Z",  
  "status": "placed",  
  "complete": true  
}
```

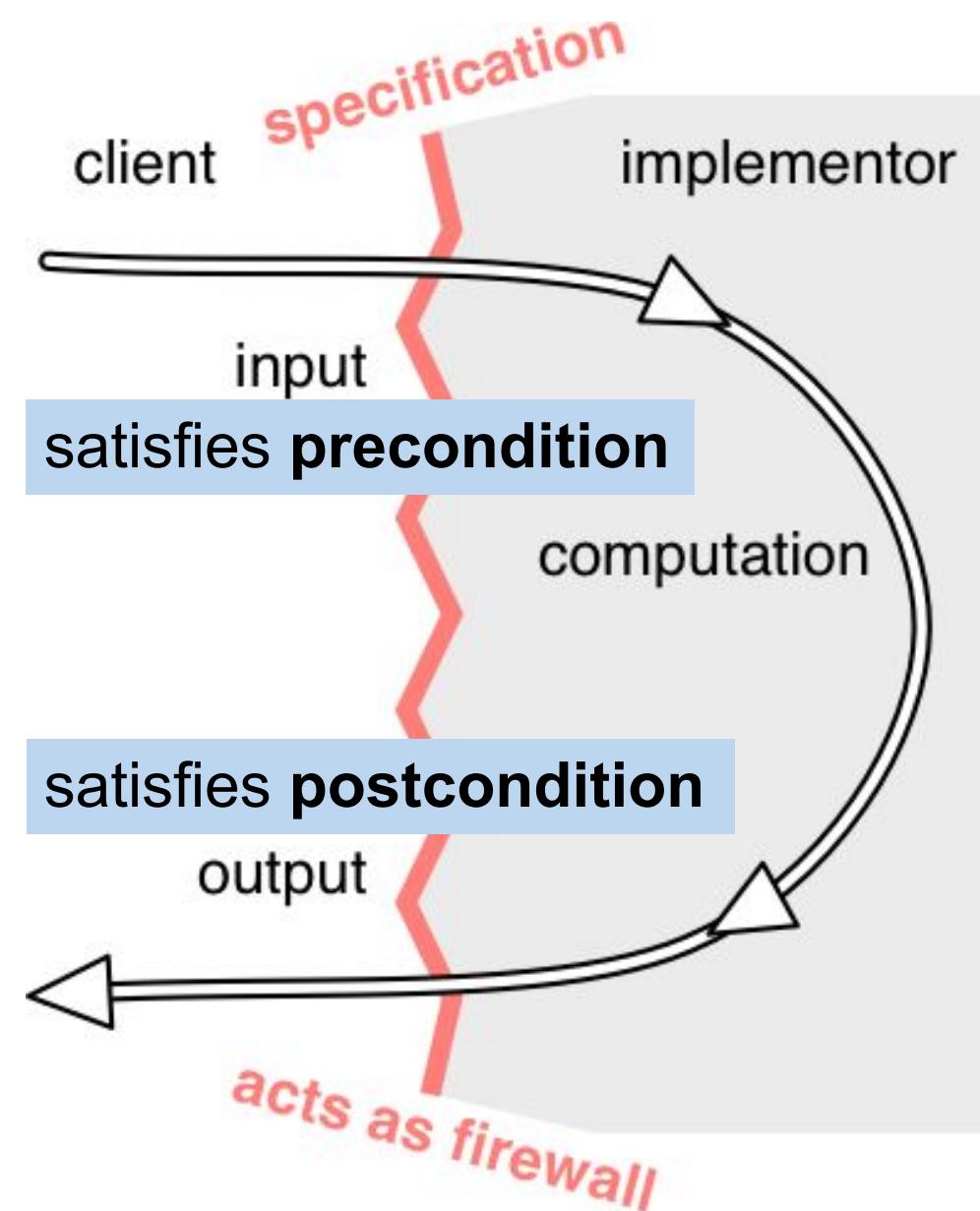
Parameter content type

application/json

9

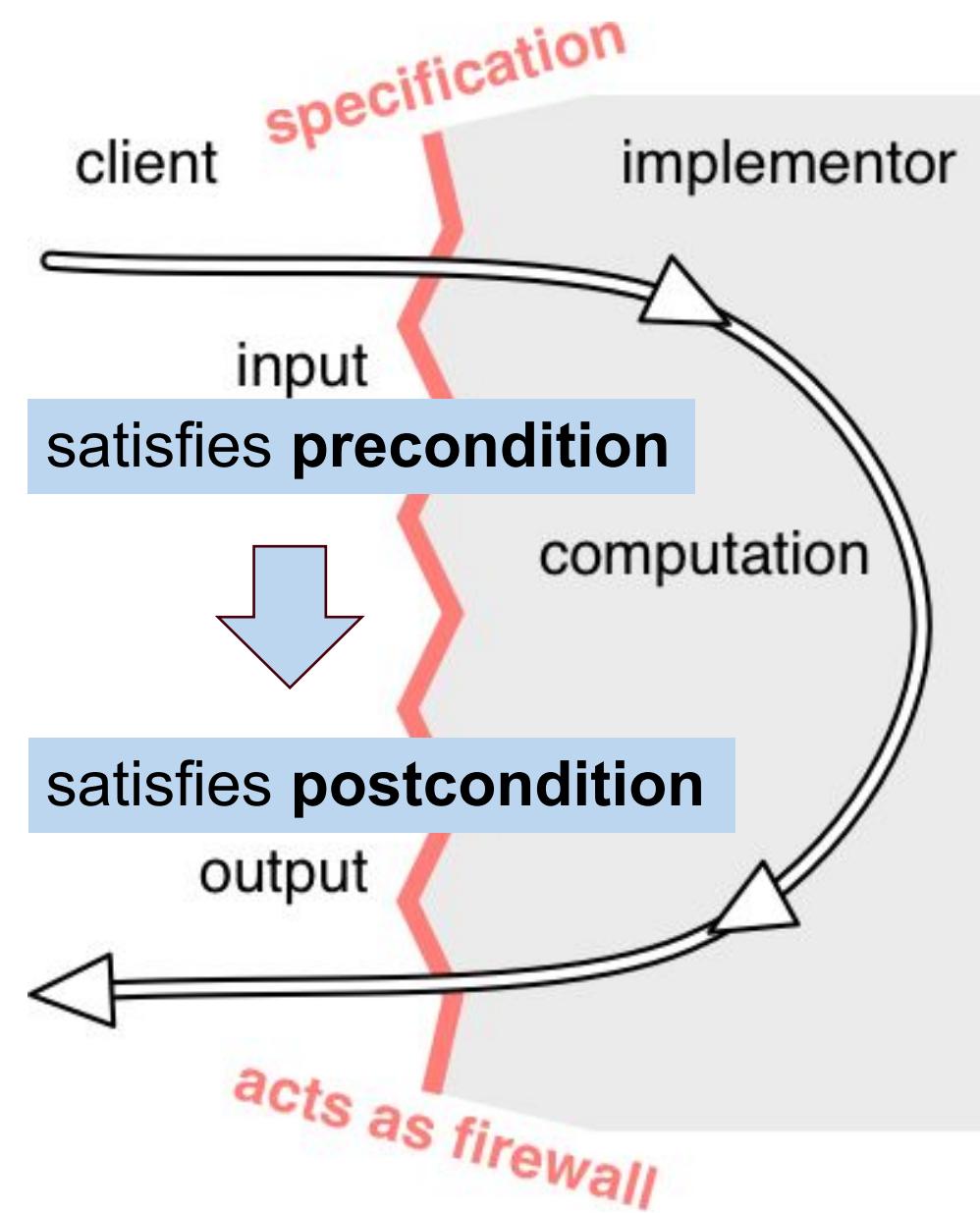
Specification: Elements

- Each specification of a function is associated with **pre- & post-conditions**
- **Pre-condition**
 - What the component **expects from the client**, expressed as a condition over the function input and/or component state
- **Post-condition**
 - What the component **promises to deliver**, as a condition over the function output and/or component state



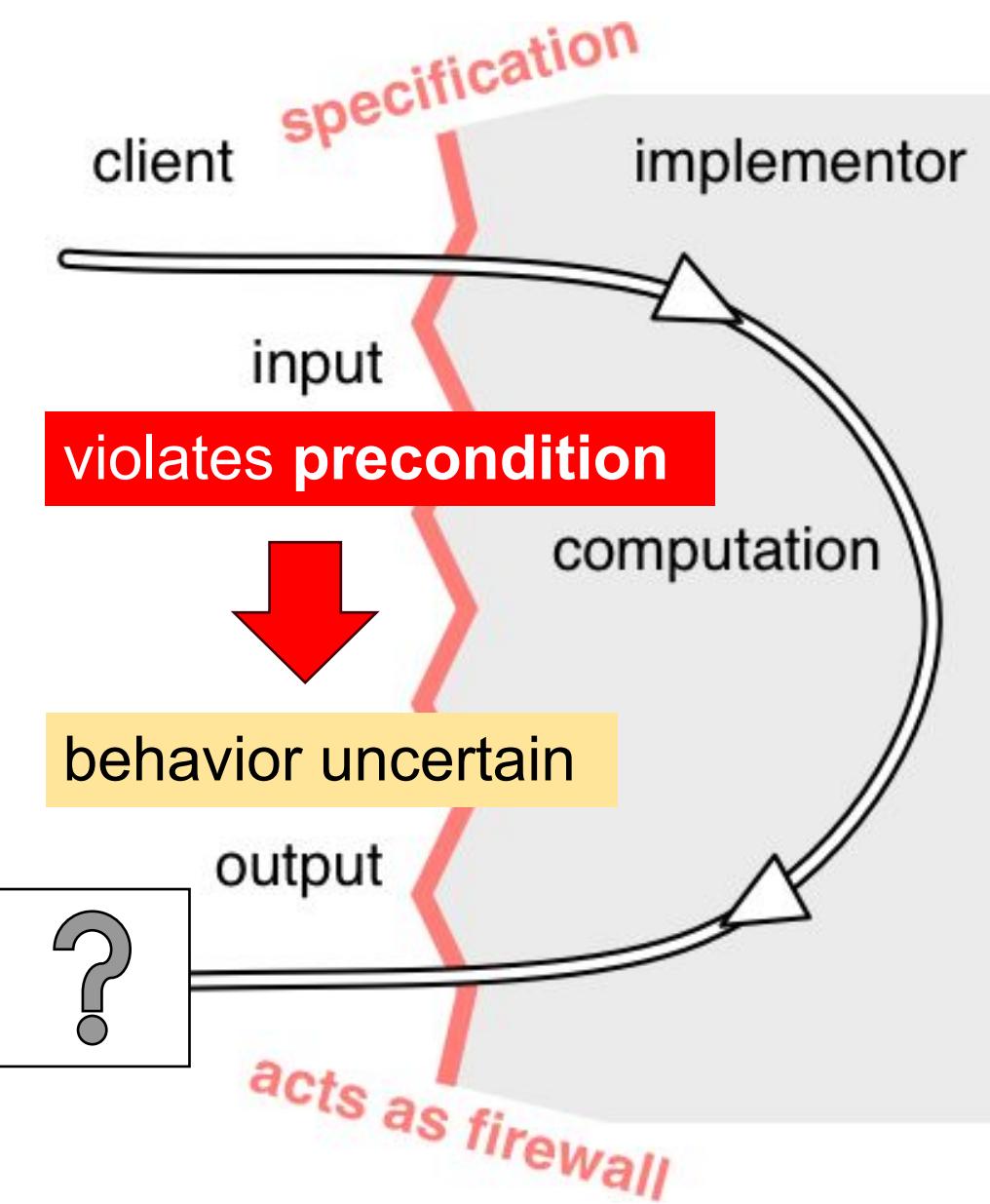
Specification: Meaning

- **Pre-condition \Rightarrow Post-condition**
(i.e., logical implication)
- If the client satisfies the pre-condition, the component promises to satisfy post-condition



Specification: Meaning

- **Pre-condition \Rightarrow Post-condition**
(i.e., logical implication)
- If the client satisfies the pre-condition, the component promises to satisfy post-condition
- But if the client violates the pre-condition, the component can behave in an arbitrary way!
 - Logically, “false implies anything”
 - **Q. Why is this reasonable?**



Example: Specifying Array Find

```
static int find(int[] arr, int val)
```

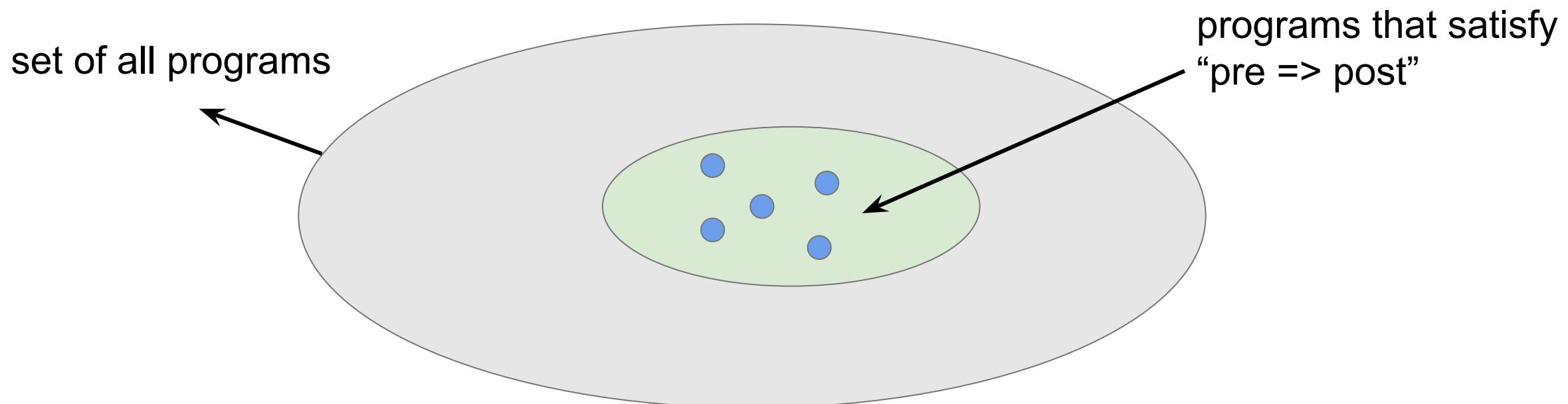
requires: val occurs exactly once in arr

effects: returns index i such that arr[i] = val

- A specification of the “find” function
- By convention, we will label pre- & post-conditions as **requires** and **effects**, respectively
- **Meaning**: If “val” occurs exactly once in “arr”, then it returns index “i” such that arr[i] = val
 - If “val” occurs zero times or more than once, then “find” may return anything

Specification as an Implementation Set

- A specification defines a **set of possible implementations**
- Given a pre- & post-condition, any implementation that fulfills the requirement “**pre-condition \Rightarrow post-condition**” is a valid implementation of the specification



Example: Implementing Array Find

```
static int find(int[] arr, int val) {  
    for (int i = 0; i < arr.length; i++) {  
        if (arr[i] == val) return i;  
    }  
    return arr.length;  
}  
  
static int find(int[] arr, int val) {  
    for (int i = arr.length - 1 ; i >= 0; i--) {  
        if (arr[i] == val) return i;  
    }  
    return -1;  
}
```

**Q. Do these functions
behave the same or
differently?**

Example: Specifying Array Find

```
static int find(int[] arr, int val)
```

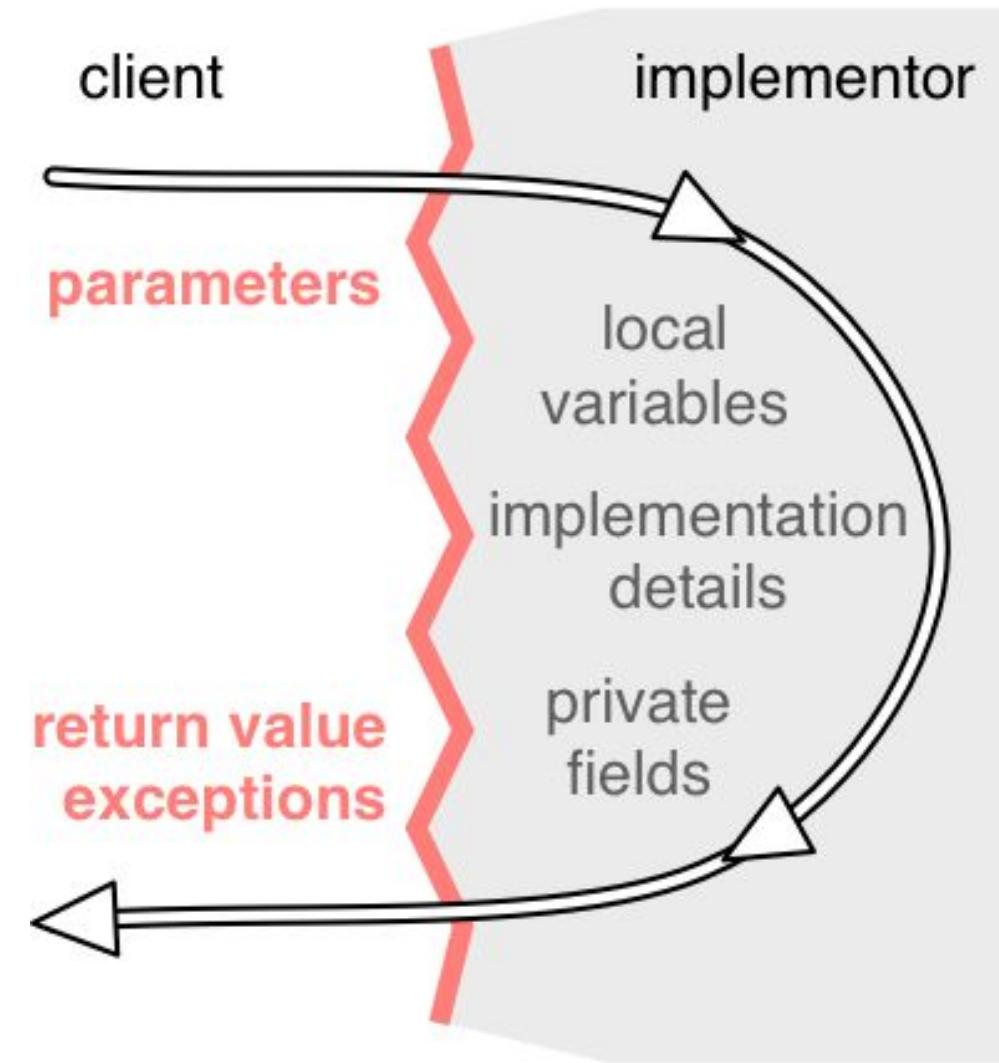
requires: val occurs exactly once in arr

effects: returns index i such that arr[i] = val

- A specification of the “find” function
- The two versions of “find” are both valid implementations of this specification!
 - As far as the client is concerned, they have the same behavior
 - One could be substituted with the other, without affecting the code of the client that relies on this specification

Specification Must Hide Unnecessary Details

- What can appear inside the pre- & post-conditions?
- Recommended practice
 - Pre-conditions should only mention input parameters of a function (**Q. Why not output?**)
 - Post-conditions should only mention the input & output parameters
 - They should avoid mentioning **hidden/private** fields in the component (**Q. Why?**)
 - If necessary, instead refer to **publicly visible** fields/functions



Specification Must Hide Unnecessary Details

```
public class Account {  
    private String accountID;  
    private int currBalance; // in cents  
  
    public void deposit(int dollars)  
        requires: nothing  
        effects: increase currBalance by (dollars)*100  
    { ... // implementation }  
}
```

- **Q. What's undesirable about this specification of “deposit”?**
- **Q. How would you improve this?**

How do we design a “good” specification?

Factors to Consider in Interface Specifications

- Deterministic vs. under-determined
- Declarative vs. operational
- Strong vs. weak
- General vs. restrictive

Deterministic vs. Under-determined

- A specification of a function is **deterministic** if, for any given input, it allows exactly one possible output.
- A specification is **under-determined** if, for some input, it allows multiple possible outputs.

Recall: Specification of Find

```
static int find(int[] arr, int val)
```

requires: val occurs exactly once in arr

effects: returns index i such that arr[i] = val

- An example of a **deterministic** specification
 - Only one return value is possible for any given input

Recall: Specification of Find

```
static int find(int[] arr, int val)
```

requires: val occurs exactly once in arr

effects: returns index i such that arr[i] = val

Spec ver1

```
static int find(int[] arr, int val)
```

requires: val occurs in arr

effects: returns index i such that arr[i] = val

Spec ver2

- **Q1.** Is the second specification (ver2) deterministic or under-determined? Why?
- **Q2.** Which of the two would you prefer? As a client? As an implementor?

Recall: Implementations of Find

```
static int find(int[] arr, int val) {  
    for (int i = 0; i < arr.length; i++) {  
        if (arr[i] == val) return i;  
    }  
    return arr.length;  
}
```

```
static int find(int[] arr, int val) {  
    for (int i = arr.length - 1 ; i >= 0; i--) {  
        if (arr[i] == val) return i;  
    }  
    return -1;  
}
```

These are both valid implementations of Spec ver1 & ver2!

Deterministic vs. Under-determined

- A specification of a function is **deterministic** if, for any given input, it allows exactly one possible output.
- A specification is **under-determined** if, for some input, it allows multiple possible outputs.
- An **under-determined** specification is ambiguous and can result in behaviors that are “surprising” to the client
 - The client can't rely on what output the function will return
- In general, **deterministic** specifications are preferable
- **Design consideration**: For a given input, are multiple outputs possible? If so, how do I modify the pre- or post-condition to make it deterministic?

Declarative vs. Operational

- An **operational** specification describes **how** a function achieves its post-condition through a series of steps
- A **declarative** specification describes **what** a function achieves without saying **how**

Declarative vs. Operational: Example

```
static int find(int[] arr, int val)
```

requires: val occurs in arr

effects: examines a[0],a[1],..., in turn and returns
the index of the 1st element equal to val

- An example of an operational specification
- **Q. What is undesirable about this specification?**

Declarative vs. Operational: Example

```
static int find(int[] arr, int val)
```

requires: val occurs in arr

effects: examines a[0],a[1],..., in turn and returns
the index of the 1st element equal to val

- An example of an operational specification
- **Q. What is undesirable about this specification?**
 - Expose details about how the function is implemented internally
 - Unnecessarily constrains the set of possible implementations

Declarative vs. Operational: Example

```
static int find(int[] arr, int val)
```

requires: val occurs in arr

effects: examines a[0],a[1],..., in turn and returns
the index of the 1st element equal to val

Operational

```
static int find(int[] arr, int val)
```

requires: val occurs in arr

effects: returns index i such that arr[i] = val

Declarative

- Declarative specifications tend to:
 - Be more concise, easier to understand
 - Allow a larger set of implementations
 - Give more flexibility to the implementor!

Declarative vs. Operational

- An **operational** specification describes **how** a function achieves its post-condition through a series of steps
- A declarative specification describes **what** a function achieves without saying **how**
- **Operational** specifications tend to:
 - Expose details about how the function is implemented internally
 - Unnecessarily constrains the set of possible implementations
- **Declarative** specifications are preferable
- **Design consideration**: Is the specification describing “how” something is done? If so, can we rewrite it to say only “what” it does?

Strong vs. Weak

- Let S1 and S2 be specifications with the same pre-condition
- S1 is **stronger** than S2 if S1 provides more guarantees about the output than S2 does
- (Mathematically, S1's post-condition is **logically stronger** than S2's post-condition)

Strong vs. Weak: Example

```
static int find(int[] a, int val)
```

requires: val occurs at least once in a

effects: returns index i such that a[i] = val

Spec ver1

```
static int find(int[] a, int val)
```

requires: val occurs at least once in a

effects: returns **lowest** index i such that a[i] = val

Spec ver2

- Spec ver2 is stronger than ver1, since it provides stronger guarantees about the output
- How strong is “strong enough”?
 - Depends on the client’s requirements
 - To fulfill their own tasks, does the client rely on the index being the lowest?

Strong vs. Weak: Example #2

```
static int find(int[] a, int val)
  requires: nothing
  effects: returns index i such that a[i] = val
```

Spec ver3

- **Q. What is wrong with ver3?**
 - The specification is too strong. In fact, there is no possible valid implementation for this specification!

Strong vs. Weak: Example #2

```
static int find(int[] a, int val)  
  requires: nothing  
  effects: returns index i such that a[i] = val
```

Spec ver3

```
static int find(int[] a, int val)  
  requires: nothing  
  effects: if val doesn't occur in a, returns -1  
        else returns index i such that a[i] = val
```

Spec ver4

- Specification should be **as weak as possible**
 - Stronger specifications allow a smaller set of implementations & are harder to implement
 - Weaker specifications give more flexibility to the implementor

Strong vs. Weak

- Let S1 and S2 be specifications with the same pre-condition
- S1 is **stronger** than S2 if S1 provides more guarantees about the output than S2 does
- A specification should be **strong enough** to support the needs of the client
- At the same time, a specification should be **as weak as possible**, to provide as flexibility to the implementor
- **Design consideration**: Is the specification providing more guarantees than needed? If so, how much can we relax them without breaking the client's code?

General vs. Restrictive

- Let S_1 and S_2 be specifications with the same post-condition
- S_1 is **more general** than S_2 if S_1 puts less restrictions on the input than S_2 does
- (Mathematically, S_1 's pre-condition is **logically weaker** than S_2 's pre-condition)

General vs. Restrictive: Example

```
static int find(int[] a, int val)
```

requires: val occurs exactly once in a

effects: returns index i such that a[i] = val

Spec ver1

```
static int find(int[] a, int val)
```

requires: val occurs in a

effects: returns index i such that a[i] = val

Spec ver2

- Spec ver2 is more general than ver1, since it accepts a larger set of inputs
 - In ver1, the client must do more work to ensure that “val” occurs exactly once
- In general, a **more general specification is preferable**, as it puts less burden on the client

General vs. Restrictive: Example #2

```
static int find(int[] a, int val)
```

requires: nothing

effects: if val doesn't occur in a, returns -1

else returns index i such that a[i] = val

Spec ver3

- Spec ver3 is most general (for the given post-condition)
- No pre-condition (“nothing”), so it accepts any inputs; no burden on the client!
- **Q. Is this always desirable?**

General vs. Restrictive: Another Example

```
static int binarySearch(int[] a, int val)
  requires: nothing
  effects: if a is not sorted or val doesn't occur in a, returns -1
            else returns index i such that a[i] = val
```

- A general specification shifts the burden onto the component to validate the input
- This can be undesirable, as it may add more complexity or performance overhead to the implementation
 - **Q. Why is this the case for the above example?**
- Sometimes, it's necessary to make the function more restrictive by strengthening the pre-condition

General vs. Restrictive

- Let S_1 and S_2 be specifications with the same post-condition
- S_1 is **more general** than S_2 if S_1 puts less restrictions on the input than S_2 does
- A specification should be **as general as possible**
 - A pre-condition places burden on the client to satisfy it
 - Less restrictive it is, more applicable the function is
- A specification should be **restrictive** when necessary
- **Design consideration**: What needs to be checked about the input for a successful operation? If the check is too expensive, can we restrict the pre-condition to rule out bad inputs?

Factors in Designing Specifications

- Deterministic vs. under-determined
- Declarative vs. operational
- Strong vs. weak
- General vs. restrictive

“Smells” for bad interface specifications

- Mentions private component details (i.e., breaking the firewall)
- Multiple possible outputs for a single input (under-determined)
- Step-by-step algorithmic descriptions (operational specs)
- Under-specified edge cases in output (too weak)
- Overly stringent output requirements (too strong)
- Overly stringent input requirements (too restrictive)
- Missing pre-condition that make post-condition impossible or inefficient to satisfy (too general)

Exercises: Are these good specifications?

static Set union(Set s1, Set s2)

requires: “s1” and “s2” are non-empty

effects: returns a new set that contains the elements from both “s1” and “s2”

static List sort(List l)

requires: nothing

effects: returns a new list that results from applying merge sort to “l”

static String read(String filepath)

requires: filepath is not null

effects: opens the file at “filepath” and returns the content of the file as a string

Interface Specifications: Takeaway

- A specification defines a contract between a component and its clients
- A specification defines a set of valid possible implementations
- A specifications should be **deterministic** rather than **under-determined**
- A specification should be **declarative** rather than **operational**
- A specification should be sufficiently **strong**, while being as **weak** as possible
- A specification should be as **general** as possible, while being **restrictive** when necessary

Summary

- Exit ticket!